

# Theories of Computation: Formative Assignment 2

To be handed in on Canvas before **Thursday 17th March, 5pm GMT**

**Exercise 1 (Time Complexity in Big-O)** Let us consider two algorithms.

Algorithm  $A_1$  has running time  $T_1(n) = \begin{cases} 5n^3 + 2 & \text{for } 0 \leq n \leq 3 \\ 7n + 9 & \text{for } n \geq 4 \end{cases}$

Algorithm  $A_2$  has running time  $T_2(n) = \begin{cases} 3n^4 + 3 & \text{for } 0 \leq n \leq 2 \\ 2n^2 & \text{for } n \geq 3 \end{cases}$

1. Show that  $T_1(n)$  is  $O(n)$  and  $T_2(n)$  is  $O(n^2)$ . [2 marks]

**Remember:** To justify your claim that  $f(n)$  is  $O(g(n))$  provide constants  $M$  and  $C$  that satisfy the property that  $\forall n \geq M. f(n) \leq Cg(n)$ .

2. For each  $n \geq 0$  which algorithm is more efficient? Justify your answer. [2 marks]

**Solution 1** 1.

$T_1(n) = O(n)$ : For  $n \geq 4$ ,  $T_1(n) = 7n + 9 \leq 10n$ . So we can take  $M_1 = 4$  and  $C_1 = 10$ .

$T_2(n) = O(n^2)$ : For  $n \geq 3$ ,  $T_2(n) = 2n^2 \leq 2n^2$ . So we can take  $M_2 = 3$  and  $C_2 = 2$ .

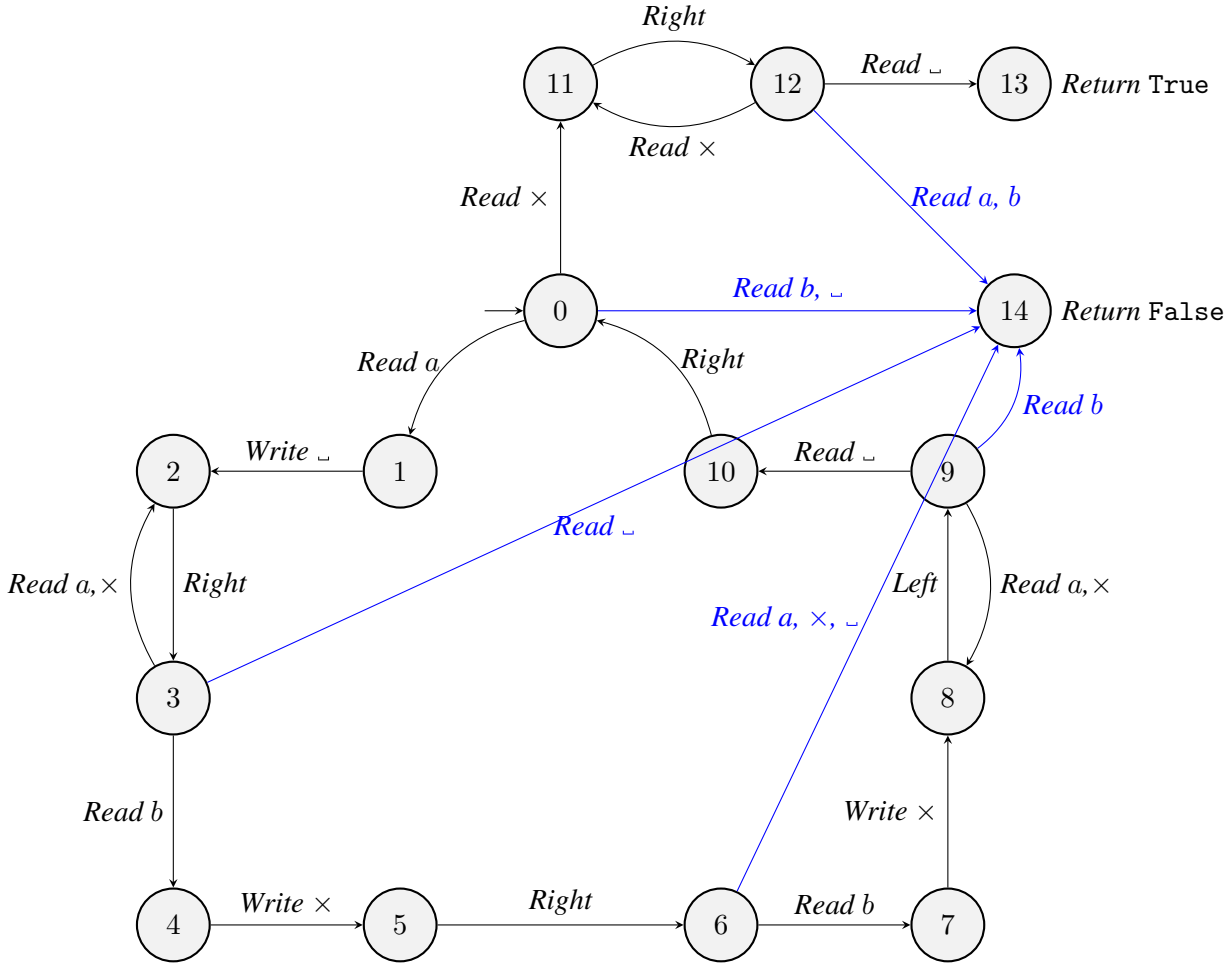
2.

<b>n</b>	0	1	2	3	4	5	6	7
<b>T<sub>1</sub>(n)</b>	2	7	42	137	37	44	51	58
<b>T<sub>2</sub>(n)</b>	3	6	51	18	32	50	72	98

For  $n = 1$ ,  $n = 3$  and  $n = 4$ ,  $A_2$  is more efficient. For  $n = 0$ ,  $n = 2$  and  $n \geq 5$ ,  $A_1$  is more efficient.

Indeed, when  $n \geq 4$ ,  $T_1(n) = T_2(n) \Leftrightarrow 2n^2 - 7n - 9 = 0 \Leftrightarrow n = 4.5$ .

**Exercise 2 (Turing Machines)** Let us consider a Turing machine  $M$  on the input alphabet  $\Sigma = \{a, b\}$  with initial state 0, tape alphabet  $T = \{a, b, \sqcup, \times\}$ , return values  $V = \{\text{True}, \text{False}\}$ , and whose transition function is represented as the following diagram.



Initially the tape contains a non-empty block of as and bs and is otherwise blank. The head is positioned on the first non blank character.

1. Give the complete run of the machine  $M$  above on the word  $ab$ . At each step, indicate the tape contents, the position of the head, the current state and the instruction (including the result if it is a Read). [2 marks]

**Hint:** No more than 10 steps are needed.

2. Without justification, does  $M$  accept words  $abb$  and  $abbb$  (i.e. return True if given them as input)? [2 marks]
3. What is the language  $\mathcal{L}(M)$  recognised by  $M$ ? [2 marks]
4. Use machine  $M$  as a macro and design a Turing machine with five states that recognises the language  $\mathcal{L} = \{a^{n+1}b^{2n} \mid n \geq 1\}$ ? [2 marks]

**Solution 2** 1.  $M$  does not accept word  $ab$ .

␣	<sup>•</sup> $a$	$b$	␣	0	Read $a$
␣	<sup>•</sup> $a$	$b$	␣	1	Write ␣
␣	␣	<sup>•</sup> $b$	␣	2	Right
␣	␣	<sup>•</sup> $b$	␣	3	Read $b$
␣	␣	<sup>•</sup> $b$	␣	4	Write $\times$
␣	␣	$\times$	␣	5	Right
␣	␣	$\times$	<sup>•</sup>	6	Read ␣
␣	␣	$\times$	<sup>•</sup>	14	Return False

2.  $M$  accepts word  $abb$  but does not accept word  $abbb$ . (Full runs below were not expected as answers.)

␣	<sup>•</sup> $a$	$b$	$b$	␣	0	Read $a$
␣	<sup>•</sup> $a$	$b$	$b$	␣	1	Write ␣
␣	␣	<sup>•</sup> $b$	$b$	␣	2	Right
␣	␣	<sup>•</sup> $b$	$b$	␣	3	Read $b$
␣	␣	<sup>•</sup> $b$	$b$	␣	4	Write $\times$
␣	␣	<sup>•</sup> $\times$	$b$	␣	5	Right
␣	␣	$\times$	<sup>•</sup> $b$	␣	6	Read $b$
␣	␣	$\times$	<sup>•</sup> $b$	␣	7	Write $\times$
␣	␣	$\times$	<sup>•</sup> $\times$	␣	8	Left
␣	␣	$\times$	$\times$	␣	9	Read $\times$
␣	␣	$\times$	$\times$	␣	8	Left
␣	␣	$\times$	$\times$	␣	9	Read ␣
␣	␣	$\times$	$\times$	␣	10	Right
␣	␣	$\times$	$\times$	␣	0	Read $\times$
␣	␣	$\times$	$\times$	␣	11	Right
␣	␣	$\times$	<sup>•</sup> $\times$	␣	12	Read $\times$
␣	␣	$\times$	<sup>•</sup> $\times$	␣	11	Right
␣	␣	$\times$	$\times$	<sup>•</sup>	12	Read ␣
␣	␣	$\times$	$\times$	<sup>•</sup> $b$	14	Return False

␣	<sup>•</sup> $a$	$b$	$b$	$b$	␣	0	Read $a$
␣	<sup>•</sup> $a$	$b$	$b$	$b$	␣	1	Write ␣
␣	␣	<sup>•</sup> $b$	$b$	$b$	␣	2	Right
␣	␣	<sup>•</sup> $b$	$b$	$b$	␣	3	Read $b$
␣	␣	<sup>•</sup> $b$	$b$	$b$	␣	4	Write $\times$
␣	␣	<sup>•</sup> $\times$	$b$	$b$	␣	5	Right
␣	␣	$\times$	<sup>•</sup> $b$	$b$	␣	6	Read $b$
␣	␣	$\times$	<sup>•</sup> $b$	$b$	␣	7	Write $\times$
␣	␣	$\times$	<sup>•</sup> $\times$	$b$	␣	8	Left
␣	␣	$\times$	$\times$	$b$	␣	9	Read $\times$
␣	␣	$\times$	$\times$	$b$	␣	8	Left
␣	␣	$\times$	$\times$	$b$	␣	9	Read ␣
␣	␣	$\times$	$\times$	$b$	␣	10	Right
␣	␣	$\times$	$\times$	$b$	␣	0	Read $\times$
␣	␣	$\times$	$\times$	$b$	␣	11	Right
␣	␣	$\times$	<sup>•</sup> $\times$	$b$	␣	12	Read $\times$
␣	␣	$\times$	<sup>•</sup> $\times$	$b$	␣	11	Right
␣	␣	$\times$	$\times$	<sup>•</sup> $b$	␣	12	Read $b$
␣	␣	$\times$	$\times$	<sup>•</sup> $b$	␣	14	Return False

3.  $\mathcal{L}(M) = \{a^n b^{2n} \mid n \geq 1\}$

4. Using machine  $M$  as a macro instruction at state 2, we can design the following machine to recognise  $\mathcal{L}$ :

