

UNIVERSITY OF BIRMINGHAM

School of Computer Science

Mathematical and Logical Foundations of Computer Science

Second Class Test 2021/22

This test is designed to be solved in about one hour and is worth 8% of your total grade.

The module was taught differently last year.
Only attempt the questions circled in red,
not the striked out ones

Mathematical and Logical Foundations of Computer Science

Question 1 [Relations, functions, induction & linear equations]

- (a) Let $V = \{1, 2, 3, 4, 5, 6, 7\}$ and let $E = \{(1, 7), (2, 1), (4, 1), (6, 5), (6, 6)\}$ be a binary relation on V . Find the equivalence closure of this relation and state the equivalence classes. (It may help to draw a diagram.) **[4 marks]**

- (b) The coefficients of the following system are taken from $\text{GF}(2)$. Solve it using Gaussian elimination.

$$\begin{aligned}x_1 + x_2 + x_4 &= 0 \\x_1 + x_3 + x_4 &= 1 \\x_2 + x_5 &= 0 \\x_1 + x_2 + x_3 + x_5 &= 0 \\x_1 + x_3 &= 0\end{aligned}$$

- (c) Let S be the smallest subset of $\{a, b\}^*$ such that all of the following conditions are satisfied:

- $\varepsilon \in S$.
- If $w \in S$, then $aabw \in S$.
- If $w \in S$, then any anagram of w is also in S . (An anagram of w is a string that arises from w by a permutation of the letters.)

- (i) Show that $aabababaa$ is in S .

- (ii) Argue that $ababab$ is not in S by giving a property and proving that all elements of S satisfy this property.

- (d) Consider the following Java methods. Do they represent functions? If yes, are they injective, surjective, or bijective? Justify your answers.

```
int doubleInt(int number) { return number * 2; }  
float addOneToFloat(float number) { return number + 1.0; }
```

Question 2 [SAT & Predicate Logic]

- (a) Let p, q, r, s be atoms capturing the states of four cells, which can either be filled or empty: p is true if the cell is filled, and false if the cell is empty, and similarly for the other atoms. Consider the following formula:

$$(p \vee \neg q) \wedge (p \vee r) \wedge (p \vee s) \wedge (q \vee \neg p) \wedge (q \vee r) \wedge (q \vee s) \\ \wedge (\neg r \vee \neg p) \wedge (\neg r \vee \neg q) \wedge (\neg r \vee s) \wedge (\neg s \vee \neg p) \wedge (\neg s \vee \neg q) \wedge (\neg s \vee r)$$

- (i) Using DPLL, find a valuation that shows that the above formula is satisfiable. Justify your answer as we did in the SAT lecture. **[6 marks]**
- (ii) Is the formula valid? Justify your answer. **[2 marks]**
- (iii) Explain in one sentence what property about the states of the four cells p, q, r , and s , this formula captures. **[2 marks]**

- (b) Consider the following signature:

- Function symbols: `zero` (arity 0); `succ` (arity 1)
- Predicate symbols: `<` (arity 2)

We will use infix notation for the binary symbol `<`. Consider the following formulas that capture properties of the above symbols:

- let S_1 be $\forall y. (\exists x. x < y) \rightarrow 0 < y$
- let S_2 be $\forall x. x < \text{succ}(x)$

For simplicity we write 0 for `zero`, 1 for `succ(zero)`, 2 for `succ(succ(zero))`, etc.

- (i) Provide a constructive Natural Deduction proof of:

$$S_1 \rightarrow S_2 \rightarrow 0 < 2$$

(Hint: you can prove this formula without $[\forall I]$ and $[\exists E]$.) **[6 marks]**

- (ii) Explain why the following tree is not a Natural Deduction proof. Justify your answer.

$$\frac{\frac{\frac{\overline{S_2}^1}{x < \text{succ}(x)} [\forall E]}{\forall y. y < \text{succ}(x)} [\forall I]}{\exists x. \forall y. y < x} [\exists I] \\ \frac{}{S_2 \rightarrow \exists x. \forall y. y < x}^1 [\rightarrow I]$$

(Hint: keep in mind that the $\forall$ and $\exists$ rules have side conditions.) **[4 marks]**