# Mathematical and Logical Foundations of Computer Science Class Test

# Solutions

Resit Class Test 2020/21

# Mathematical and Logical Foundations of Computer Science Class Test

# Question 1 [Math]

### (a) Set Theory

(i) Consider the binary relation R on  $\mathbb{Z}_5$  given by

$$R = \{(x, y) \in \mathbb{Z}_5 \times \mathbb{Z}_5 \mid x + y = 0\}$$

List the elements of R.

#### [2 marks]

Is *R* reflexive, symmetric, transitive? Is it an equivalence relation? Give brief justifications or counterexamples. [4 marks]

(ii) Discuss the similarities and differences between Z and the elements of the Java datatype int. [4 marks]

#### (b) Linear Algebra

Consider the following points in 3D space

$$P_1 = \begin{pmatrix} 1 \\ 0 \\ 2 \end{pmatrix} \quad P_2 = \begin{pmatrix} 0 \\ 1 \\ 1 \end{pmatrix} \quad P_3 = \begin{pmatrix} -1 \\ -3 \\ 2 \end{pmatrix}$$

- (i) Show that they form a *scalene* triangle (i.e., a triangle in which no two sides have the same length). [3 marks]
- (ii) The triangle defines a plane E in  $\mathbb{R}^3$ . Give its parametric representation and its normal form. [3 marks]
- (iii) Find the nearest neighbour of  $Q = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}$  on *E*. [2 marks]

(iv) Find a basis for  $\mathbb{R}^3$  containing two sides of the triangle. [2 marks]

#### Model answer / LOs / Creativity:

- (a) (i)  $R = \{(0, 0), (1, 4), (2, 3), (3, 2), (4, 1)\}$ . R is
  - not reflexive because  $(1, 1) \notin R$ ;
  - symmetric because x + y = y + x in  $\mathbb{Z}_5$ ;
  - not transitive because (2, 3) and (3, 2) belong to R but (2, 2) does not;
  - not an equivalence relation because it is neither reflexive nor transitive.

- (ii) Some properties discussed in the course:
  - $\mathbb{Z}$  has infinitely many elements, int has  $2^{32}$  many.
  - $\bullet$  On both  ${\mathbb Z}$  and int addition, subtraction, and multiplication are defined.
  - Both  ${\mathbb Z}$  and int satisfy the ring laws.
  - $\bullet~\mathbb{Z}$  satisfies the cancellation law for multiplication, int does not.

(Each of these, if properly stated, may attract two points.)

(b) (i) We compute the length of each side:

$$|\overrightarrow{P_1P_2}| = \sqrt{(0-1)^2 + (1-0)^2 + (1-2)^2} = \sqrt{1+1+1} = \sqrt{3}$$
$$|\overrightarrow{P_2P_3}| = \sqrt{(-1-0)^2 + (-3-1)^2 + (2-1)^2} = \sqrt{1+16+1} = \sqrt{18} = 3\sqrt{2}$$
$$|\overrightarrow{P_3P_1}| = \sqrt{(1-(-1))^2 + (0-(-3))^2 + (2-2)^2} = \sqrt{4+9+0} = \sqrt{13}$$

and find that they are all different.

# (ii) One parametric representation is $X = \begin{pmatrix} 1 \\ 0 \\ 2 \end{pmatrix} + s \cdot \begin{pmatrix} -1 \\ 1 \\ -1 \end{pmatrix} + t \cdot \begin{pmatrix} -1 \\ -4 \\ 1 \end{pmatrix}$ , using

a corner and two sides of the triangle.

A normal to the plane, according to the formula in the handout is given by  $\vec{n} = \begin{pmatrix} -3 \\ 2 \\ 5 \end{pmatrix}$ . Using a corner of the triangle we get the normal form  $-3x_1 + 2x_2 + 5x_3 = 7$ .

(iii) Using the formula in the handout we compute the nearest neighbour as  $Q' = \begin{pmatrix} -3 \end{pmatrix}$ 

$$Q + \frac{7}{38} \cdot \begin{pmatrix} -3\\2\\5 \end{pmatrix} = \frac{7}{38} \cdot \begin{pmatrix} -3\\2\\5 \end{pmatrix}.$$

(iv) We can choose two sides of the triangle plus the normal:  $B = \{ \begin{pmatrix} -1 \\ 1 \\ -1 \end{pmatrix}, \begin{pmatrix} -1 \\ -4 \\ 1 \end{pmatrix}, \begin{pmatrix} -3 \\ 2 \\ 5 \end{pmatrix} \}.$ 

(Because the normal is at a right angle to the triangle, it is linearly independent of the sides of the triangle.)

Learning outcomes: "Solve mathematical problems in algebra and set theory" (a & b); "Apply mathematical techniques to solve a problem within a computer science setting" (a-i).

## Question 2 [Logic]

#### (a) **Propositional logic**

(i) Let *F* be the following proposition:  $(\neg(P \lor Q)) \rightarrow ((\neg P) \rightarrow (\neg Q) \rightarrow R) \rightarrow R$ . Provide an intuitionistic Natural Deduction proof of *F*. **[5 marks]** 

[1 mark]

- (ii) Is F valid? Justify your answer.
- (iii) Consider the formula F, where  $\lor$  is turned into an  $\land$ , i.e., let G be  $(\neg(P \land Q)) \rightarrow$  $((\neg P) \rightarrow (\neg Q) \rightarrow R) \rightarrow R$ . Is G valid? Is it satisfiable? Justify your answers. [4 marks]

#### (b) **Predicate Logic**

Consider the following signature:

- Function symbols: zero (arity 0); succ (arity 1)
- Predicate symbols: < (arity 2)

We will use infix notation for the binary symbol <. Consider the following formulas that capture properties of the above symbols:

- let  $S_1$  be  $\forall x.x < \operatorname{succ}(x)$
- let  $S_2$  be  $\forall x. \neg (x < x)$
- (i) Provide a constructive Sequent Calculus proof of:

$$S_1, S_2 \vdash \neg (\forall x. \forall y. x < y \rightarrow \operatorname{succ}(x) < y)$$

You are allowed to write 0 for zero, 1 for succ(zero), 2 for succ(succ(zero)), etc. (Hint: 9 proof steps are enough to complete this proof, but feel free to use a different number of proof steps in your proof.) [6 marks]

(ii) Provide a model  $M_1$  such that  $\models_{M_1} \neg (\forall x.\forall y.x < y \rightarrow \text{succ}(x) < y)$ , and a model  $M_2$  such that  $\models_{M_2} \forall x.\forall y.x < y \rightarrow \text{succ}(x) < y$  [4 marks]

#### Model answer / LOs / Creativity:

(a) (i) Here is an intuitionistic Natural Deduction proof of F:

$$\frac{\overline{\neg (P \lor Q)} \stackrel{1}{=} \frac{\overline{P} \stackrel{3}{\rightarrow} [\lor I_{L}]}{[\neg E]}}{\frac{\neg P \rightarrow \neg Q \rightarrow R}{\stackrel{2}{=} \frac{1}{\neg P} \stackrel{3}{[\neg P]} [\rightarrow E]} \stackrel{[\lor I_{L}]}{[\rightarrow E]} \frac{\overline{\neg (P \lor Q)} \stackrel{1}{=} \frac{\overline{Q} \stackrel{4}{\rightarrow} [\lor I_{R}]}{[\neg P \lor Q} [\neg E]} \frac{\frac{1}{\neg Q} \stackrel{4}{[\neg I]} [\neg E]}{[\rightarrow E]} \frac{\frac{1}{\neg Q} \stackrel{4}{[\neg I]} [\neg E]}{[\rightarrow E]} \frac{\overline{((\neg P) \rightarrow (\neg Q) \rightarrow R) \rightarrow R} \stackrel{2}{\rightarrow R} \stackrel{2}{[\rightarrow I]} [\rightarrow I]}{[\rightarrow I]}$$

(ii) By consistency of Natural Deduction, because F is provable it is also valid.

(iii) G is not valid because if P is true and both Q and R are false, then G is also false. It is however satisfiable because if P, Q, and R are all false, then G is true.

$$\begin{array}{c|c} \hline \hline \hline \hline S_1, S_2 \vdash 0 < 1 \\ \hline \hline S_1, S_2 \vdash 0 < 1 \\ \hline \hline S_1, S_2 \vdash 0 < 1 \\ \hline \hline S_1, S_2 \vdash 0 < 1 \\ \hline \hline S_1, S_2, 1 < 1 \vdash \bot \\ \hline \hline S_1, S_2, 1 < 1 \vdash \bot \\ \hline \hline \hline S_1, S_2, 0 < 1 \rightarrow 1 < 1 \vdash \bot \\ \hline \hline \hline S_1, S_2, \forall y.0 < y \rightarrow 1 < y \vdash \bot \\ \hline \hline \hline S_1, S_2, \forall y.0 < y \rightarrow 1 < y \vdash \bot \\ \hline \hline \hline S_1, S_2, \forall x. \forall y. x < y \rightarrow \operatorname{succ}(x) < y \vdash \bot \\ \hline \hline \hline \hline S_1, S_2 \vdash \neg (\forall x. \forall y. x < y \rightarrow \operatorname{succ}(x) < y) \\ \hline \end{array}$$

(ii) The model  $M_1 = \langle \mathbb{N}, \langle 0, +1 \rangle, \langle \{ \langle n, m \rangle \mid n < m \} \rangle \rangle$ , where +1 is the function that given a number increments it by 1, is a model of  $\neg (\forall x.\forall y.x < y \rightarrow \operatorname{succ}(x) < y)$ , and the model  $M_2 = \langle \mathbb{N}, \langle 0, +1 \rangle, \langle \{ \langle n, m \rangle \mid \operatorname{True} \} \rangle \rangle$  is a model of  $\forall x.\forall y.x < y \rightarrow \operatorname{succ}(x) < y$ 

Learning outcomes: "Understand and apply algorithms for key problems in logic such as satisfiability." (a); "Write formal proofs for propositional and predicate logic" (a & b); "Apply logical techniques to solve a problem within a computer science setting" (a & b).